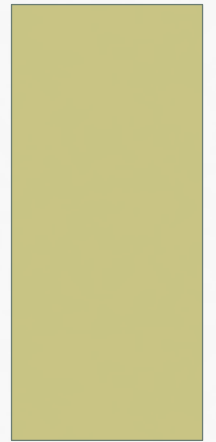


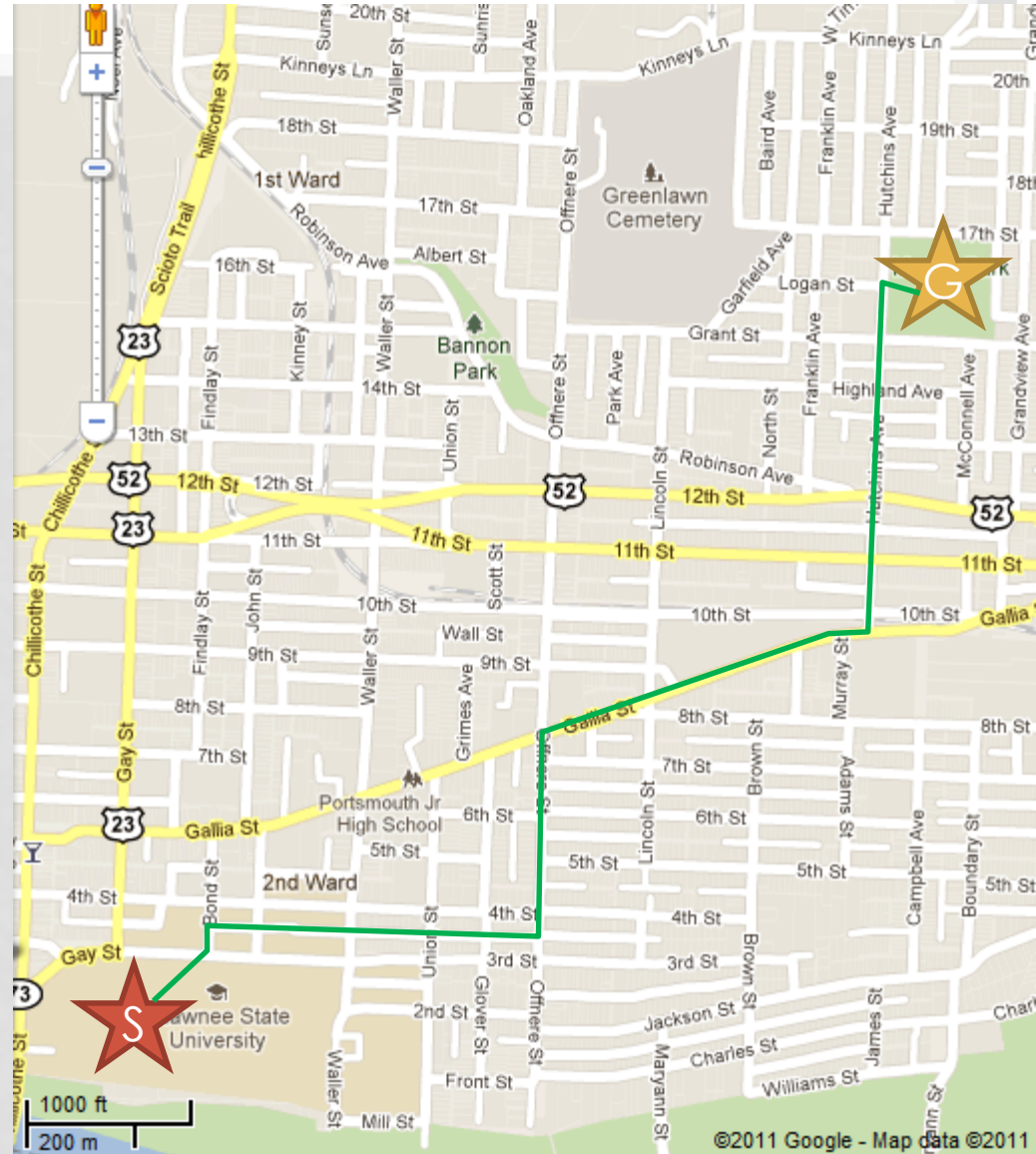
PATHFINDING

REFERENCE: "ARTIFICIAL INTELLIGENCE FOR GAMES", IAN MILLINGTON.



GOALS

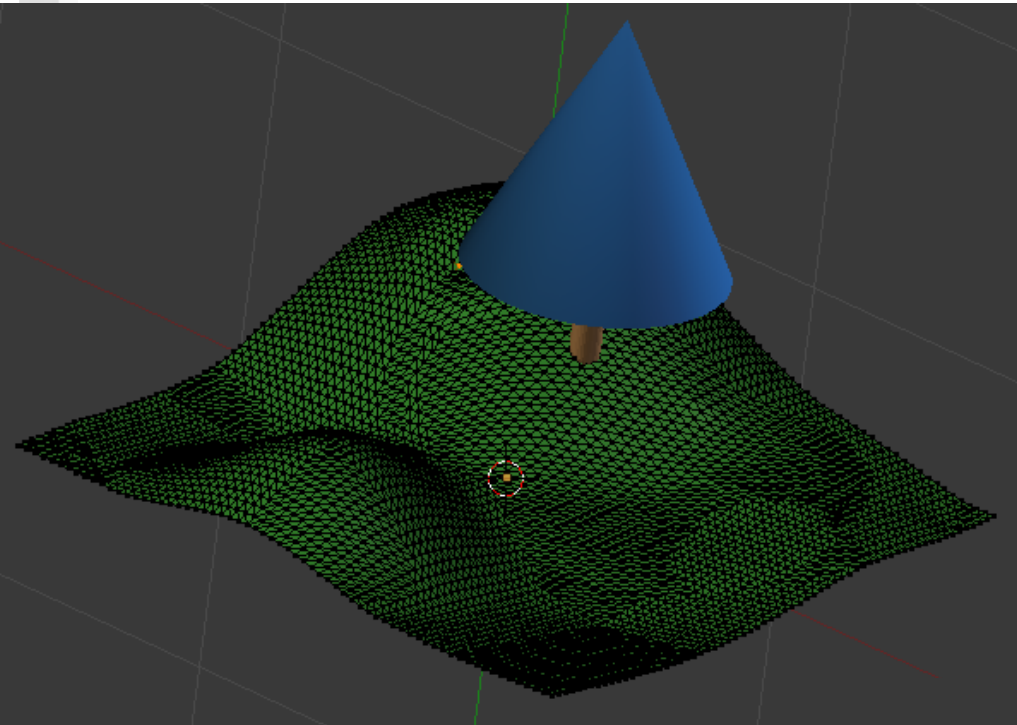
- Find a “pretty-short” path from point A to point B.
 - The *best* path is often prohibitively expensive to find.



NAV-MESHES

- A common way to represent:
 - Walkable areas of a map.
 - Cover-spots
 - Jumping-points (to cross a chasm)
 - Ladders
 - Save points, ammo drops, etc.
- Key Ideas:
 - Not usually visible
 - Usually much lower poly-count than the actual ground
- We'll use these in this lab to represent the **search graph**.

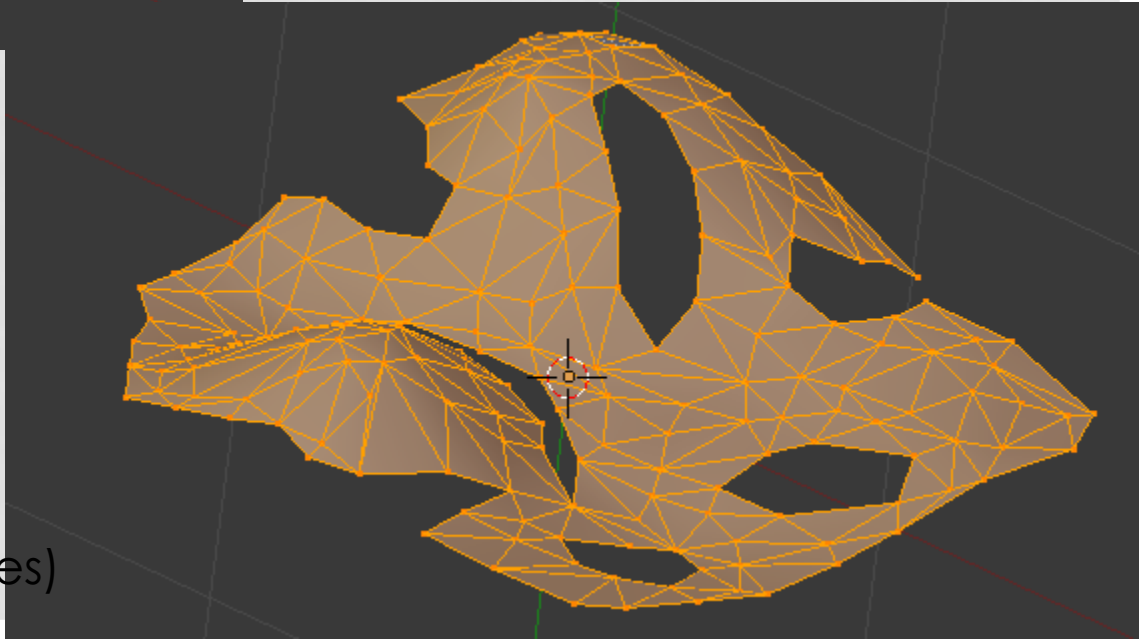
NAV-MESHES, CONT.



Display Mesh (16541 faces)



Nav Mesh (301 faces)



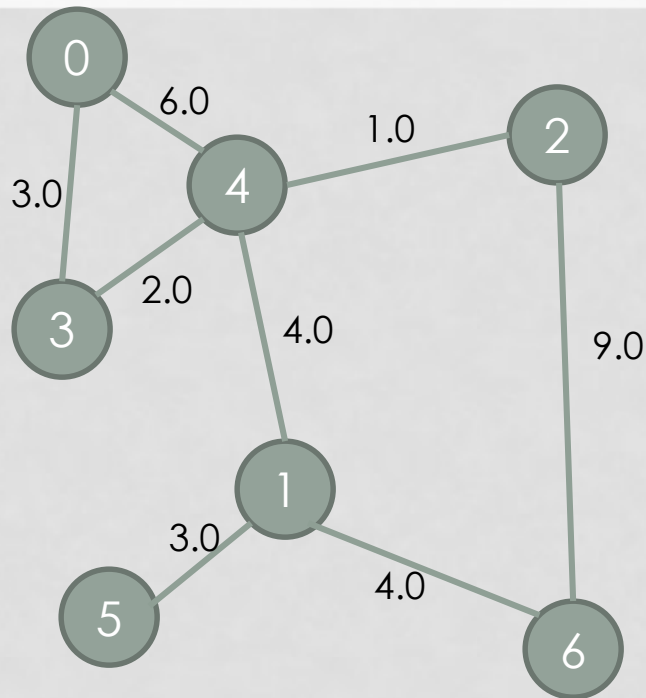
NAV-MESHES, CONT.

- We can use these for pathfinding...
 - Nodes are faces (their center?)
 - Edges are connections between neighboring faces
 - Cost is just the euclidean distance between centers.
- Probably best done off-line
 - Finding neighbors can be costly.

GRAPH-BASED ALGORITHM

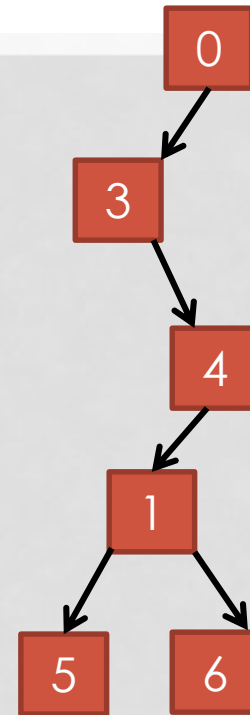
- Graph
 - Vertices (Nodes)
 - Edges
 - Single edge
 - Double edge
 - Cost
 - (For map traversal), distance is the norm
 - Cost multiplier (>1) for rough terrain.
- Representation
 - Adjacency Matrix
 - Adjacency List
- [Connections to NavMesh]

GRAPH NODES VS. SEARCH NODES



Graph

Start == 0, Goal == 6



Search Tree

There are usually many search trees for the same graph.

GRAPH NODES VS. SEARCH NODES, CONT.

- Recommended approach:
 - GraphNode class:
 - [Any map-related data]
 - Neighbor_**pointers** = [...]
 - SearchNode class (I actually just used a tuple here):
 - GraphNode reference
 - [optional] Heuristic value (later in A*)
 - Cost-so-far value
 - SearchTree map
 - Key = GraphNode reference
 - Value = parent GraphNode reference (or null for root)
 - I actually stored this plus the depth and heuristic.
- Create SearchNodes and add to map as you explore the map.
 - Positive: Easy to tell if you've visited a node – it's in the map.
 - Positive: Easy to restart the search – just throw out the map
 - DON'T make any references to the search in the GraphNodes!

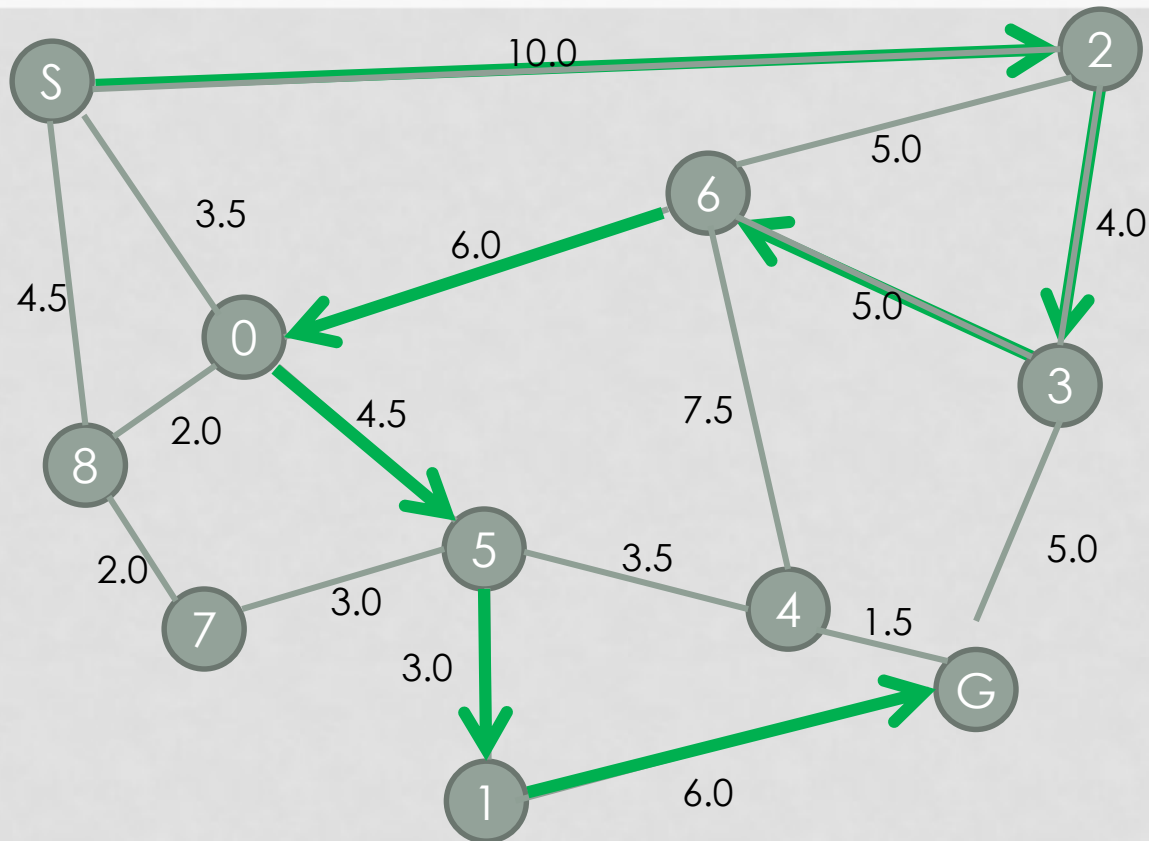
ACTUAL SEARCH

- Start with initial node, add it to the map
- Which node do we visit next?
 - Depth-First Search
 - Breadth-First Search
 - Dijkstra Search
 - A* Search
 - ...

DEPTH-FIRST P.C.

```
# Could be faster w/ a stack (non-recursive)
def dfs(curSNode, parentMap, mainGraph, g2sMap):
    for n in curSNode.GNode.neighbors:
        # [optional] early termination if n is goal.
        if n not in g2sMap:
            newCost = curSNode.costSoFar +
                mainGraph.edgeCost(curSNode.GNode, n)
            newSNode = new SNode(n, newCost)
            parentMap[newSNode] = curSNode
            g2sMap[n] = newSNode
            dfs(newSNode, parentMap, mainGraph, g2sMap)
initSNode = new SNode(start, 0)
searchTree = {initSNode : null}
graph2searchMap = {start : initSNode}
mapGraph = ...
dfs(initSNode, searchTree, mapGraph, graph2searchMap)
```

DEPTH-FIRST SEARCH



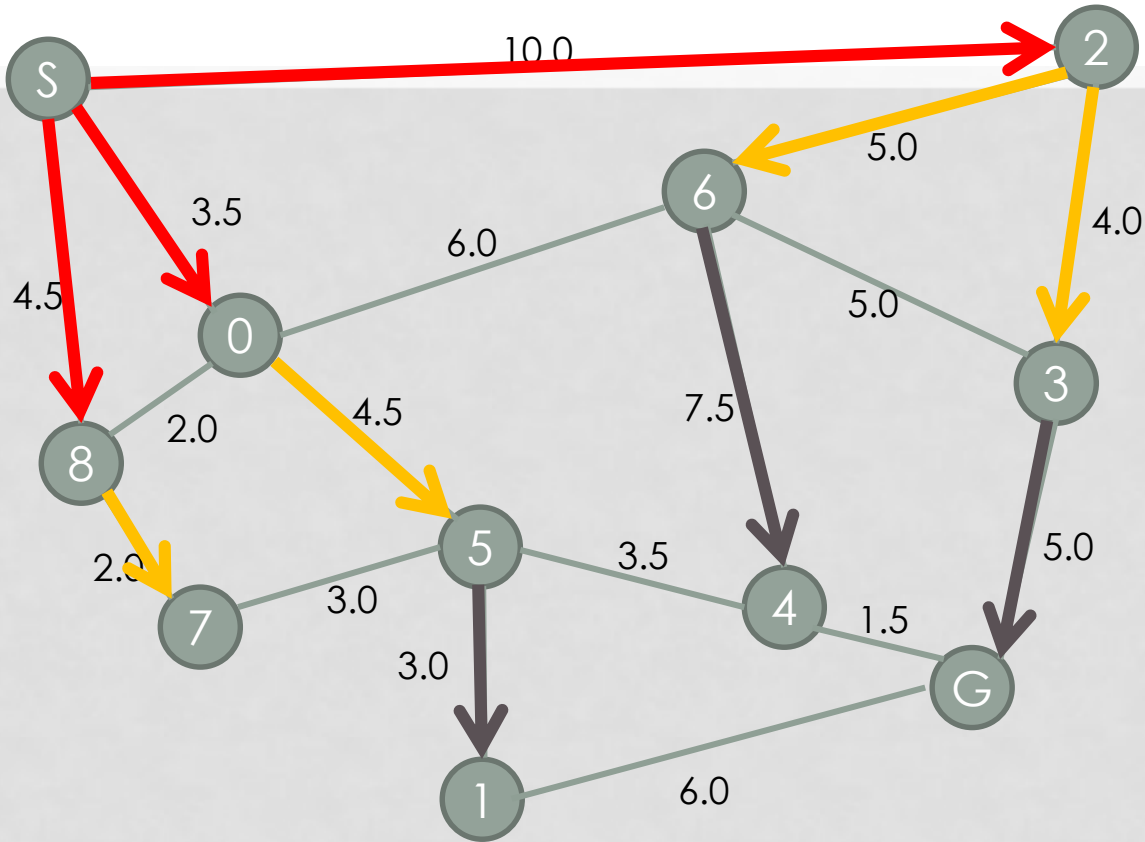
Total=38.5

The path we get will, in general, not be the best

BREADTH-FIRST SEARCH P.C.

```
initSNode = new SNode(start, 0)
searchTree = {initSNode : null}
graph2searchMap = {start : initSNode}
mapGraph = ...
frontier = [initSNode]
while len(frontier) != 0:
    new_frontier = []
    for cur in frontier:
        # [optional] early termination if
        # cur is the goal
        for n in cur.GNode.neighbors:
            if n not in graph2searchMap:
                newCost = cur.costSoFar +
                    mapGraph.edgeCost(s.GNode, n)
                newSNode = new SNode(n, newCost)
                searchTree[newSNode] = cur
                graph2searchMap[n] = newSNode
                new_frontier.append(newSNode)
    frontier = new_frontier
```

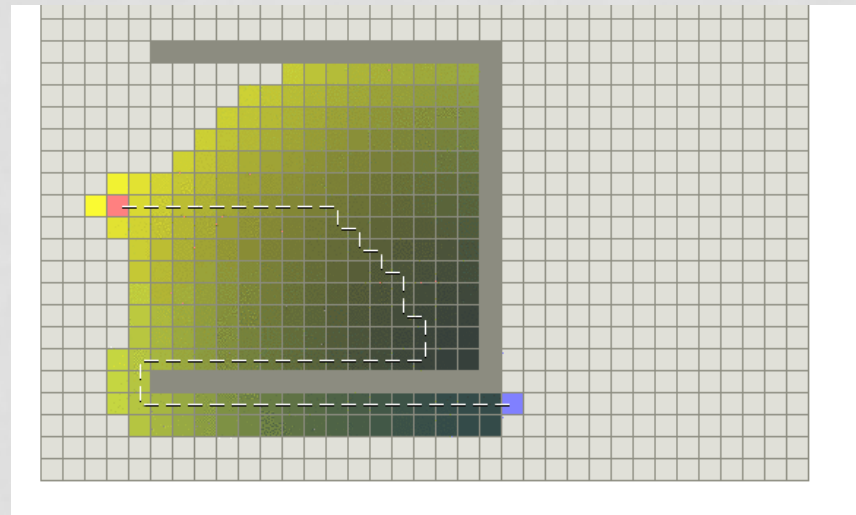
BREADTH-FIRST SEARCH



Again, we don't necessarily get the "best" path.

GREEDY-FIRST

- Just like BFS, but pick the node that *appears* to take us closer to goal
 - For this, we need some way to calculate a **heuristic**: a guess at distance to goal
 - This is usually an under-estimate – see why?
- Problem:



DJIKSTRA'S ALGORITHM P.C.

- Like BFS also, but solves the greedy-first problem.
- But...does more work than necessary.

```
# ... everything else as with BFS
```

```
for n in cur.GNode.neighbors:
```

```
    newCost = cur.costSoFar +
```

```
        mapGraph.edgeCost(s.GNode, n)
```

```
    if n not in graph2searchMap:
```

```
        newSNode = new SNode(n, newCost)
```

```
        searchTree[newSNode] = cur
```

```
        graph2searchMap[n] = newSNode
```

```
    elif newCost < graph2searchMap[n].costSoFar:
```

```
        searchTree[graph2searchMap[n]] = cur
```

```
        graph2searchMap[n].costSoFar = newCost
```

- Better than BFS because we *re-route parent pointers* if we **find a better path**.
- We aren't guaranteed the absolute optimal solution, but it's a pretty-good soln.

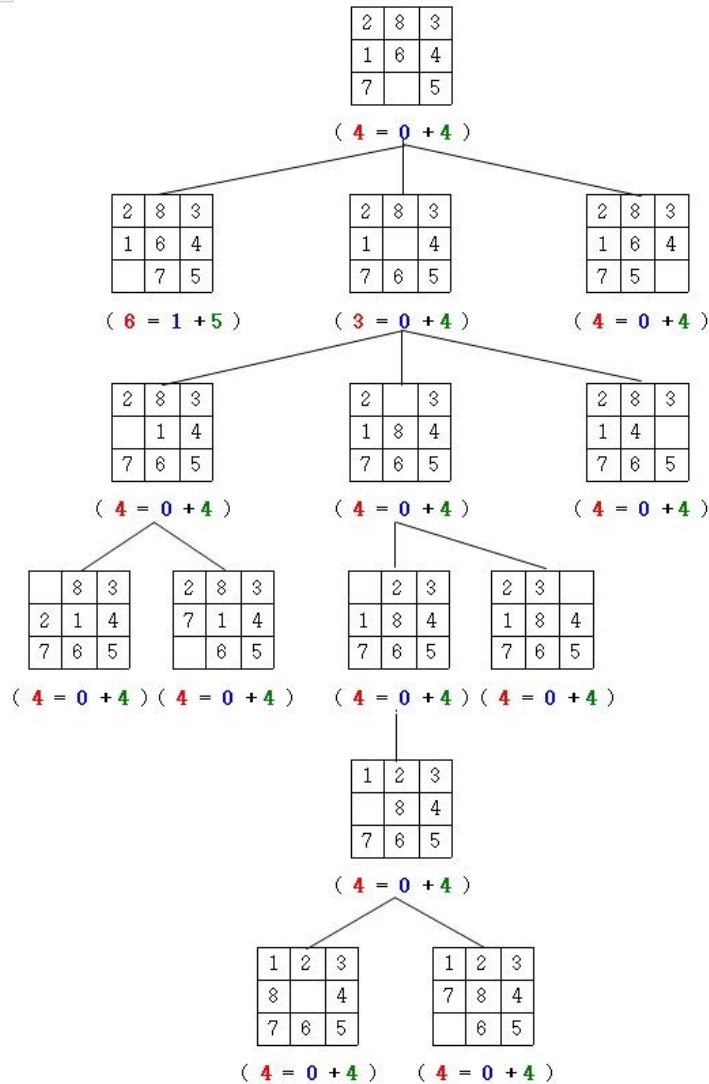
A* ALGORITHM

- Rather than looking at all neighbors in turn, pick the one we *think* will lie on the best path.
 - A combination of actual_cost + heuristic cost
 - Note how greedy-first only considered heuristic cost.
 - $f(x) = g(x) + h(x)$ # Total = actual + heuristic
 - It is vital that the heuristic be **consistent**: i.e. less than or equal to the actual cost
 - Otherwise, nodes we've visited must be re-examined.
 - We also don't want the heuristic to be too low: otherwise we'll look at too many nodes.
- Re-route visited nodes, like Dijkstra

A* P.C.

```
initSNode = new SNode(start, 0)
searchTree = {initSNode : null}
open = new MinHeap() # Of SearchNode's. Order by totalCost (actual + heuristic)
open.push(initSNode)
closed = new set()
graph2searchMap = {start : initSNode}
mapGraph = ...
while !open.empty():
    cur = open.pop()
    closed.add(cur)
    # [not optional here] Early termination if cur is the goal
    for n in cur.GNode.neighbors:
        newCost = cur.costSoFar + mapGraph.edgeCost(s.GNode, n)
        if n not in graph2searchMap:
            newSNode = new SNode(n, newCost)
            searchTree[newSNode] = cur
            graph2searchMap[n] = newSNode
            open.push(newSNode)
        elif n not in closed and newCost < graph2searchMap[n].costSoFar:
            # Better path through us - like Djikstra's
            graph2searchNode[n].costSoFar = newCost
            searchTree[graph2searchNode[n]] = cur
            # Important: reheapify open!
        # Not necessary if your heuristic is consistent. If not, we need to re-open
        # closed nodes
        # elif newCost < graph2searchMap[n].costSoFar:
        #     graph2searchNode[n].costSoFar = newCost
        #     searchTree[graph2searchNode[n]] = cur
        #     closed.remove(graph2searchNode[n])
        #     open.add(graph2searchNode[n])
return failure
```

ANOTHER APPLICATION OF A*



THE “NEW” A* ?

- Jump-Point Search seems to be gaining popularity in game programming.
- Good reference:
 - <https://harablog.wordpress.com/2011/09/07/jump-point-search/>
- Makes many improvements, but only works on square-grids.
 - Can we make this work with nav-meshes?